Introduction: Fibonacci Numbers II

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Algorithmic Design and Techniques
Algorithms and Data Structures at edX

Learning Objectives

- Produce a simple algorithm to compute Fibonacci numbers.
- Show that this algorithm is very slow.

Definition

$$F_n = \begin{cases} 0, & n = 0\,, \\ 1, & n = 1\,, \\ F_{n-1} + F_{n-2}, & n > 1\,. \end{cases}$$

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Grows rapidly.

Computing Fibonacci numbers

Compute F_n

Input: An integer $n \geq 0$.

Output: F_n.

Algorithm

FibRecurs(n)

```
if n ≤ 1:
  return n
```

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else:
  return FibRecurs(n − 1) + FibRecurs(n − 2)
```

Running time

Let $T(n)$ denote the number of lines of code executed by `FibRecurs`(n).

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$T(n) = 2.$

If $n \geq 2$

FibRecurs(n)

```
if n ≤ 1:
  return n
else:
  return FibRecurs(n − 1) + FibRecurs(n − 2)
```

If $n \geq 2$

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FibRecurs(n)

if n ≤ 1:
  return n
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$T(n) = 3$

If $n \geq 2$

FibRecurs(n)

```
if n ≤ 1:
  return n
else:
  return FibRecurs(n − 1) + FibRecurs(n − 2)
```

$$T(n) = 3 + T(n-1) + T(n-2).$$

Running Time

$$T(n) = \begin{cases} 2 & \text{if } n \leq 1 \\ T(n-1) + T(n-2) + 3 & \text{else.} \end{cases}$$

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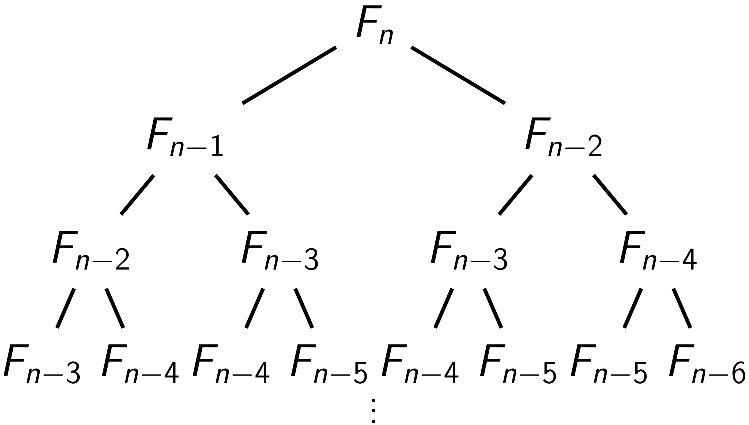
Takes 56,000 years at 1GHz.

Why so slow?



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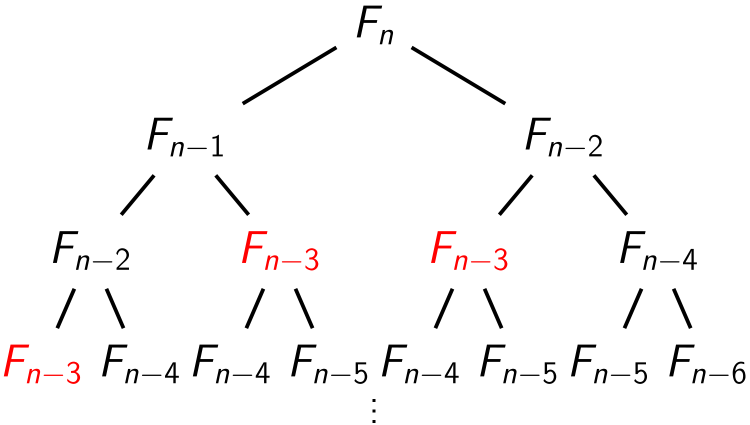

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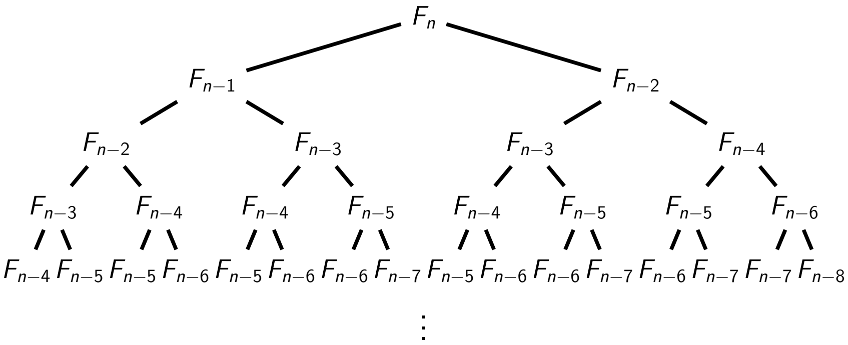
F_n
F_{n-1} F_{n-2}
F_{n-2} F_{n-3} F_{n-3} F_{n-4}
F_{n-3} F_{n-4} F_{n-4} F_{n-5} F_{n-4} F_{n-5} F_{n-5} F_{n-6}
F_{n-4} F_{n-5} F_{n-5} F_{n-6} F_{n-5} F_{n-6} F_{n-6} F_{n-7} F_{n-5} F_{n-6} F_{n-6} F_{n-7} F_{n-6} F_{n-7} F_{n-7} F_{n-8}
⋮

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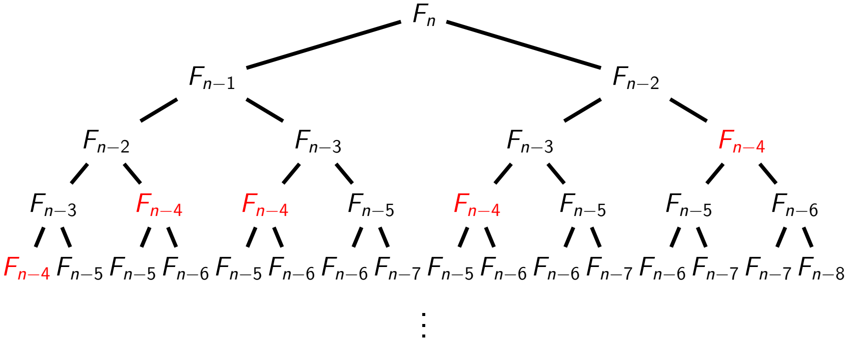
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