# Introduction: Fibonacci Numbers III

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Algorithmic Design and Techniques
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#### Learning Objectives

Compute Fibonacci numbers efficiently.

#### Definition

$$F_n = \begin{cases} 0, & n = 0, \\ 1, & n = 1, \\ F_{n-1} + F_{n-2}, & n > 1. \end{cases}$$

# Algorithm

#### FibRecurs(n)

```
\begin{aligned} &\text{if } n \leq 1\colon \\ &\text{return } n \\ &\text{else:} \\ &\text{return } \text{FibRecurs}(n-1) + \text{FibRecurs}(n-2) \end{aligned}
```

Too slow!

Imitate hand computation:

0, 1

Imitate hand computation:

0, 1, 1

0 + 1 = 1

#### Imitate hand computation:

0, 1, 1, 2

$$0 + 1 = 1$$

$$1 + 1 = 2$$

Imitate hand computation:

0, 1, 1, 2, 3

$$0 + 1 = 1$$

$$1 + 1 = 2$$

$$1 + 2 = 3$$

Imitate hand computation:

$$0 + 1 = 1$$

$$1 + 1 = 2$$

$$1 + 2 = 3$$

$$2 + 3 = 5$$

Imitate hand computation:

0, 1, 1, 2, 3, 5, 8

$$0 + 1 = 1$$

$$1 + 1 = 2$$

$$1 + 2 = 3$$

$$2 + 3 = 5$$

$$3 + 5 = 8$$

# New Algorithm

FibList(n)create an array F[0...n]

 $F[i] \leftarrow F[i-1] + F[i-2]$ 

 $F[0] \leftarrow 0$ 

 $F[1] \leftarrow 1$ 

return F[n]

for i from 2 to n:

## New Algorithm

# FibList(n)

create an array F[0...n]  $F[0] \leftarrow 0$   $F[1] \leftarrow 1$ for i from 2 to n:  $F[i] \leftarrow F[i-1] + F[i-2]$ 

- T(n) = 2n + 2. So T(100) = 202.
- Easy to compute.

return F[n]

#### Summary

- Introduced Fibonacci numbers.
- Naive algorithm takes ridiculously long time on small examples.
- Improved algorithm incredibly fast.

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Moral: The right algorithm makes all the difference.