Divide-and-Conquer: Master Theorem

Neil Rhodes

Department of Computer Science and Engineering
University of California, San Diego

Algorithmic Design and Techniques
Algorithms and Data Structures at edX

Outline

1. What is the Master Theorem

2. Proof of Master Theorem

$$T(n) = T\left(\frac{n}{2}\right) + O(1)$$

$$T(n) = T\left(\frac{n}{2}\right) + O(1)$$

$$\downarrow$$

$$T(n) = O(\log n)$$

$$T(n) = 4T\left(\frac{n}{2}\right) + O(n)$$

$$T(n) = 4T\left(\frac{n}{2}\right) + O(n)$$

$$\Big\downarrow$$

$$T(n) = O(n^2)$$

$$T(n) = 3T\left(\frac{n}{2}\right) + O(n)$$

$$T(n) = 3T\left(\frac{n}{2}\right) + O(n)$$

$$\downarrow$$

$$T(n) = O(n^{\log_2 3})$$

$$T(n) = 2T\left(\frac{n}{2}\right) + O(n)$$

$$T(n) = 2T\left(\frac{n}{2}\right) + O(n)$$

$$\Big\downarrow$$

$$T(n) = O(n \log n)$$

Master Theorem

Theorem

If $T(n) = aT\left(\left\lceil \frac{n}{b} \right\rceil\right) + O(n^d)$

Master Theorem

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If $T(n) = aT\left(\left\lceil \frac{n}{b} \right\rceil\right) + O(n^d)$ (for constants $a > 0, b > 1, d \geq 0$), then:

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If $T(n) = aT\left(\left\lceil \frac{n}{b} \right\rceil\right) + O(n^d)$ (for constants $a > 0, b > 1, d \geq 0$), then:

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Master Theorem Example 1

$$T(n) = 4T\left(\frac{n}{2}\right) + O(n)$$

Master Theorem Example 1

$$T(n) = 4T\left(\frac{n}{2}\right) + O(n)$$

$$a = 4$$

Master Theorem Example 1

$$T(n) = 4T\left(\frac{n}{2}\right) + O(n)$$

$$a = 4$$

$$b = 2$$

Master Theorem Example 1

$$T(n) = 4T\left(\frac{n}{2}\right) + O(n^{1})$$

$$a = 4$$

$$b = 2$$

$$d = 1$$

Master Theorem Example 1

$$\begin{aligned} T(n) &= 4T\left(\frac{n}{2}\right) + O(n) \\ a &= 4 \\ b &= 2 \\ d &= 1 \end{aligned}$$

Since $d < \log_b a$, $T(n) = O(n^{\log_b a}) = O(n^2)$

Master Theorem Example 2

$$T(n) = 3T\left(\frac{n}{2}\right) + O(n)$$

Master Theorem Example 2

$$T(n) = 3T\left(\frac{n}{2}\right) + O(n)$$
$$a = 3$$

Master Theorem Example 2

$$T(n) = 3T\left(\frac{n}{2}\right) + O(n)$$

$$a = 3$$

$$b = 2$$

Master Theorem Example 2

$$T(n) = 3T\left(\frac{n}{2}\right) + O(n^{1})$$

$$a = 3$$

$$b = 2$$

$$d = 1$$

Master Theorem Example 2

$$T(n) = 3T\left(\frac{n}{2}\right) + O(n)$$
$$a = 3$$
$$b = 2$$
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Since $d < \log_b a$,
$T(n) = O(n^{\log_b a}) = O(n^{\log_2 3})$

Master Theorem Example 3

$$T(n) = 2T\left(\frac{n}{2}\right) + O(n)$$

Master Theorem Example 3

$$T(n) = 2T\left(\frac{n}{2}\right) + O(n)$$

$$a = 2$$

Master Theorem Example 3

$$T(n) = 2T\left(\frac{n}{2}\right) + O(n)$$

$$a = 2$$

$$b = 2$$

Master Theorem Example 3

$$T(n) = 2T\left(\frac{n}{2}\right) + O(n^{1})$$

$$a = 2$$

$$b = 2$$

$$d = 1$$

Master Theorem Example 3

$$T(n) = 2T\left(\frac{n}{2}\right) + O(n)$$

$$a = 2$$

$$b = 2$$

$$d = 1$$

Since $d = \log_b a$,
$T(n) = O(n^d \log n) = O(n \log n)$

Master Theorem Example 4

$$T(n) = T\left(\frac{n}{2}\right) + O(1)$$

Master Theorem Example 4

$$T(n) = 1T\left(\frac{n}{2}\right) + O(1)$$

$$a = 1$$

Master Theorem Example 4

$$T(n) = T\left(\frac{n}{2}\right) + O(1)$$

$$a = 1$$

$$b = 2$$

Master Theorem Example 4

$$T(n) = T\left(\frac{n}{2}\right) + O(n^{0})$$

$$a = 1$$

$$b = 2$$

$$d = 0$$

Master Theorem Example 4

$$T(n) = T\left(\frac{n}{2}\right) + O(1)$$

$$a = 1$$

$$b = 2$$

$$d = 0$$

Since $d = \log_b a$, $T(n) = O(n^d \log n) = O(n^0 \log n) = O(\log n)$

Master Theorem Example 5

$$T(n) = 2T\left(\frac{n}{2}\right) + O(n^2)$$

Master Theorem Example 5

$$T(n) = 2T\left(\frac{n}{2}\right) + O(n^2)$$

$$a = 2$$

Master Theorem Example 5

$$T(n) = 2T\left(\frac{n}{2}\right) + O(n^2)$$

$$a = 2$$

$$b = 2$$

Master Theorem Example 5

$$T(n) = 2T\left(\frac{n}{2}\right) + O(n^2)$$

$$a = 2$$

$$b = 2$$

$$d = 2$$

Master Theorem Example 5

$$T(n) = 2T\left(\frac{n}{2}\right) + O(n^2)$$

$$a = 2$$

$$b = 2$$

$$d = 2$$

Since $d > \log_b a$, $T(n) = O(n^d) = O(n^2)$

Outline

Master Theorem

Theorem

If $T(n) = aT\left(\left\lceil \frac{n}{b} \right\rceil\right) + O(n^d)$ (for constants $a > 0, b > 1, d \geq 0$), then:

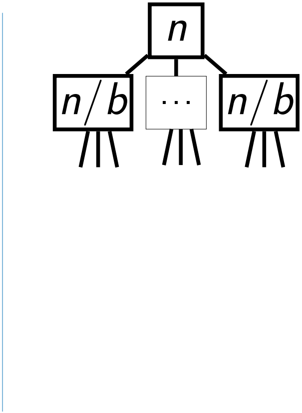
$$T(n) = \begin{cases} O(n^d) & \text{if } d > \log_b a \\ O(n^d \log n) & \text{if } d = \log_b a \\ O(n^{\log_b a}) & \text{if } d < \log_b a \end{cases}$$

$$T(n) = aT\left(\left\lceil \frac{n}{b} \right\rceil\right) + O(n^d)$$

$$\boxed{n}$$

$$T(n) = aT\left(\left\lceil \frac{n}{b} \right\rceil\right) + O(n^d)$$

level

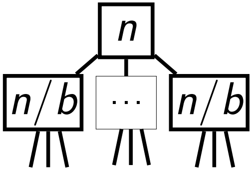


$$T(n) = aT\left(\left\lceil \frac{n}{b} \right\rceil\right) + O(n^d)$$

level

0

1



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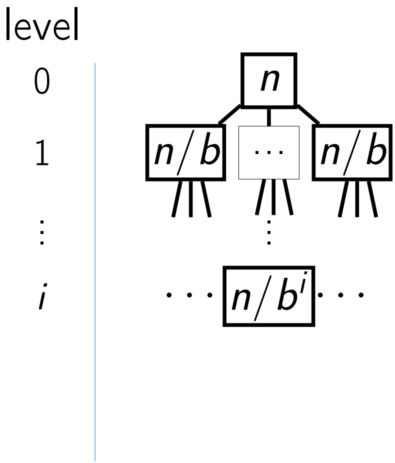

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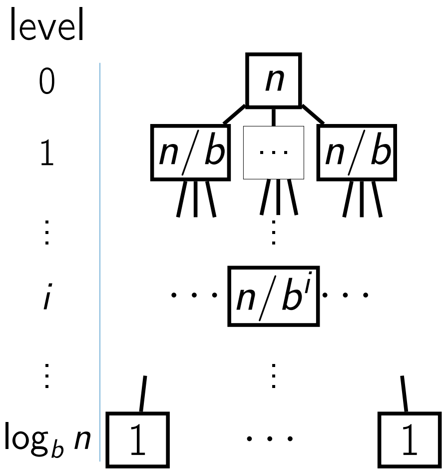

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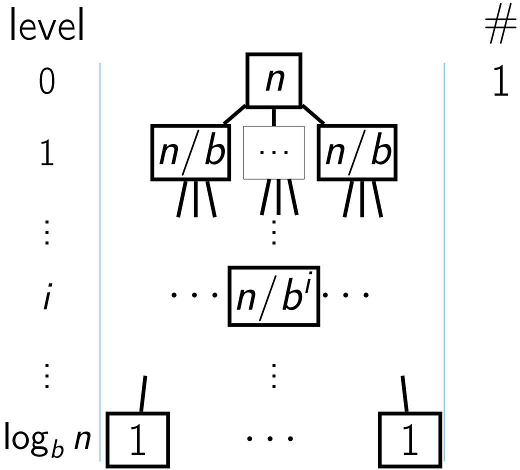

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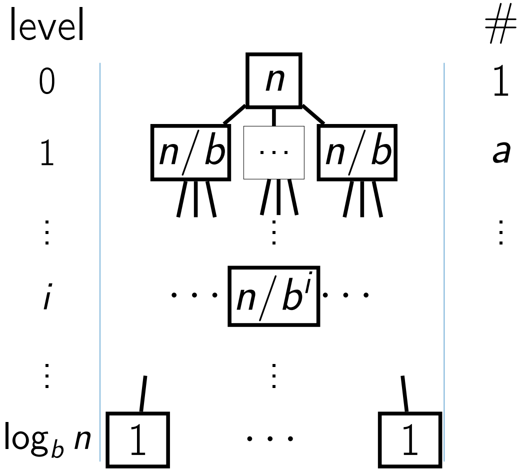

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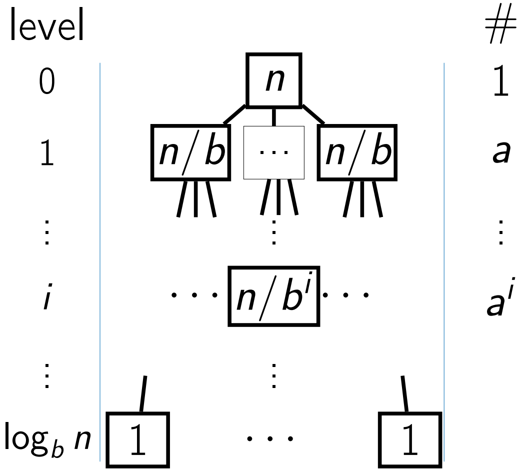

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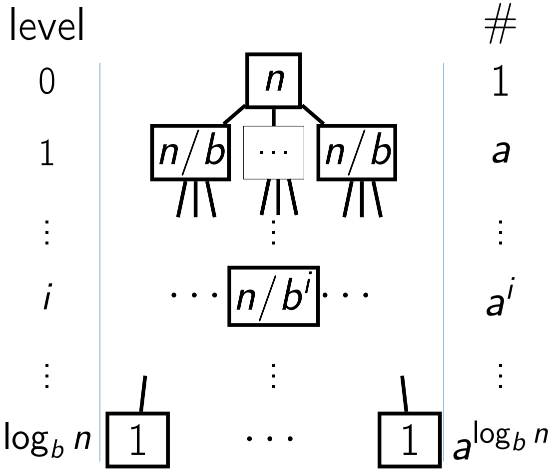

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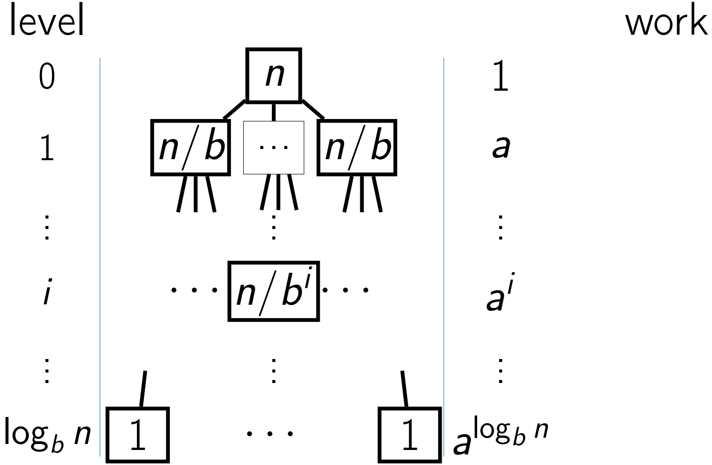

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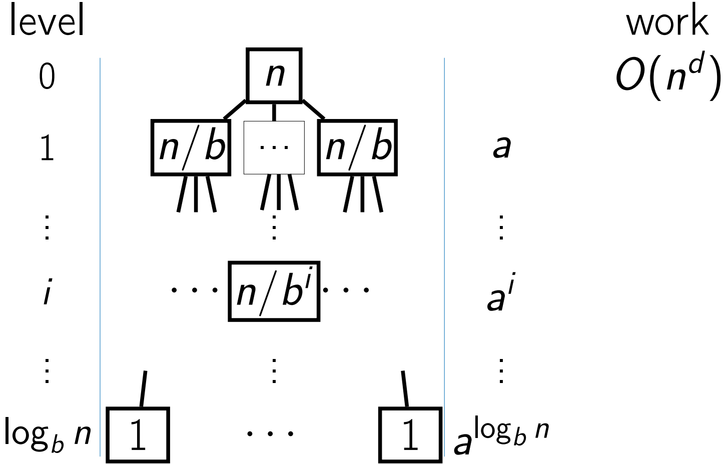

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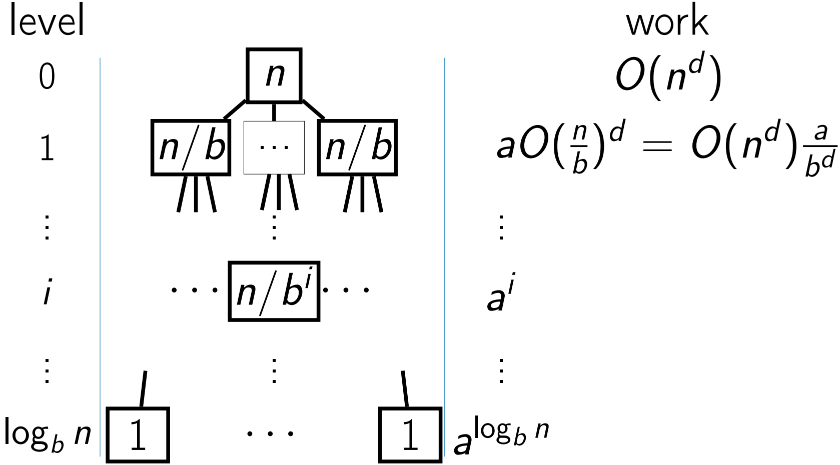

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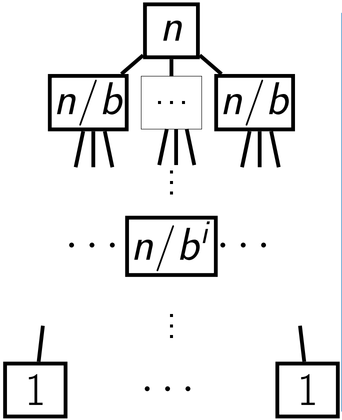
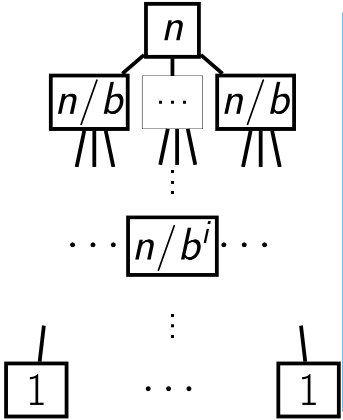


level	work
0	$O(n^d)$
1	$aO(\frac{n}{b})^d = O(n^d)\frac{a}{b^d}$
$\vdots$	$\vdots$
i	$a^i O\left(\frac{n}{b^i}\right)^d = O(n^d)\left(\frac{a}{b^d}\right)^i$
$\vdots$	$\vdots$
$\log_b n$	$a^{\log_b n}$

$$T(n) = aT\left(\left\lceil \frac{n}{b} \right\rceil\right) + O(n^d)$$



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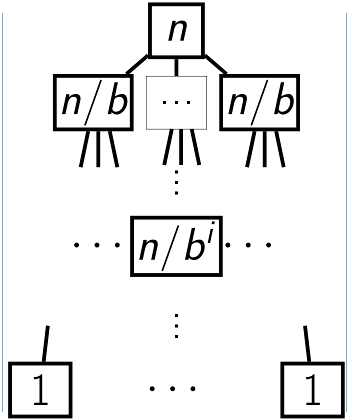
$\vdots$	$\vdots$
$\log_b n$	$a^{\log_b n} = O(n^{\log_b a})$

$$T(n) = aT\left(\left\lceil \frac{n}{b} \right\rceil\right) + O(n^d)$$



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$\vdots$	$\vdots$
i	$a^i O\left(\frac{n}{b^i}\right)^d = O(n^d)\left(\frac{a}{b^d}\right)^i$
$\vdots$	$\vdots$
$\log_b n$	$a^{\log_b n} = O(n^{\log_b a})$

Total: $\sum_{i=0}^{\log_b n} O(n^d)\left(\frac{a}{b^d}\right)^i$

Geometric Series

For $r \neq 1$:

$$a + ar + ar^2 + ar^3 + \cdots + ar^{n-1}$$

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For $r \neq 1$:

$$\begin{aligned} & a + ar + ar^2 + ar^3 + \cdots + ar^{n-1} \\ = & a\frac{1 - r^n}{1 - r} \end{aligned}$$

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$$
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& a + ar + ar^2 + ar^3 + \cdots + ar^{n-1} \\
=& a\frac{1 - r^n}{1 - r} \\
=& \left\{ \right.
\end{aligned}
$$

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For $r \neq 1$:

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\begin{aligned}
& a + ar + ar^2 + ar^3 + \cdots + ar^{n-1} \\
=& a\frac{1 - r^n}{1 - r} \\
=& \begin{cases} O(a) & \text{if } r < 1 \\ & \end{cases}
\end{aligned}
$$

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\begin{aligned}
& a + ar + ar^2 + ar^3 + \cdots + ar^{n-1} \\
= & a\frac{1 - r^n}{1 - r} \\
= & \begin{cases} O(a) & \text{if } r < 1 \\ O(ar^{n-1}) & \text{if } r > 1 \end{cases}
\end{aligned}
$$

Case 1: $\frac{a}{b^d} < 1$ $(d > log_b a)$

$$\sum_{i=0}^{\log_b n} O(n^d)\left(\frac{a}{b^d}\right)^i$$

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$$\sum_{i=0}^{\log_b n} O(n^d)\left(\frac{a}{b^d}\right)^i$$

$$= O(n^d)$$

Case 2: $\frac{a}{b^d} = 1$ ($d = log_b a$)

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$$= \sum_{i=0}^{\log_b n} O(n^d)$$

Case 2: $\frac{a}{b^d} = 1$ ($d = log_b a$)

$$
\begin{aligned}
&\sum_{i=0}^{\log_b n} O(n^d)\left(\frac{a}{b^d}\right)^i \\
=&\sum_{i=0}^{\log_b n} O(n^d) \\
=&(1 + \log_b n) O(n^d)
\end{aligned}
$$

Case 2: $\frac{a}{b^d} = 1$ ($d = log_b a$)

$$\begin{aligned}
&\sum_{i=0}^{\log_b n} O(n^d)\left(\frac{a}{b^d}\right)^i \\
=&\sum_{i=0}^{\log_b n} O(n^d) \\
=&(1+\log_b n)O(n^d) \\
=&O(n^d \log n)
\end{aligned}$$

Case 3: $\frac{a}{b^d} > 1$ ($d < log_b a$)

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$$= O\left(O(n^d)\left(\frac{a}{b^d}\right)^{\log_b n} \right)$$

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$$= O\left(O(n^d)\frac{a^{\log_b n}}{b^{d \log_b n}}\right)$$

Case 3: $\frac{a}{b^d} > 1$ ($d < log_b a$)

$$\sum_{i=0}^{\log_b n} O(n^d)\left(\frac{a}{b^d}\right)^i$$

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Case 3: $\frac{a}{b^d} > 1$ ($d < log_b a$)

$$
\begin{aligned}
&\sum_{i=0}^{\log_b n} O(n^d)\left(\frac{a}{b^d}\right)^i \\
=&O\left(O(n^d)\left(\frac{a}{b^d}\right)^{\log_b n}\right) \\
=&O\left(O(n^d)\frac{a^{\log_b n}}{b^{d\log_b n}}\right) \\
=&O\left(O(n^d)\frac{n^{\log_b a}}{n^d}\right) \\
=&O(n^{\log_b a})
\end{aligned}
$$

Summary

Master theorem is a shortcut:

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Theorem

If $T(n) = aT\left(\left\lceil \frac{n}{b} \right\rceil\right) + O(n^d)$ (for constants $a > 0, b > 1, d \geq 0$), then:

$$T(n) = \begin{cases} O(n^d) & \text{if } d > \log_b a \\ O(n^d \log n) & \text{if } d = \log_b a \\ O(n^{\log_b a}) & \text{if } d < \log_b a \end{cases}$$